

A comparative survey of methods based on Banach's contraction principle to analyze the cost of algorithms with two recurrence equations

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- 2 Product of fuzzy quasi-metric defined on the domain of words
- 3 Product quasi-metric space of complexity space
- 4 Preordered complete fuzzy quasi-metric space
- 5 Summary

Introduction

Complexity of Algorithms. 2-Recurrences Algorithm

Some authors have applied the Banach contraction principle equipped with an appropriate complexity space introduced by Schellekens, see [Sch95], [GarRomSch08] or [RomSchTiVal07], and bicomplete quasi-metrics [RoSaTi07], [RomTi09], [SaadVaCho09] to show the existence and uniqueness of solution for the recurrence equations of several well known algorithms.

Introduction

Complexity of Algorithms. 2-Recurrences Algorithm

To complement such works we have chosen a kind of algorithm composed by **two recurrences**, from [Atk96].

```
function P(n)
  if n > 0 then
    Q(n-1); C; P(n-1); C; Q(n-1)
function Q(n)
  if n > 0 then
    P(n-1); C; Q(n-1); C; P(n-1); C; Q(n-1)
```

Introduction

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```

The execution times $S(n)$ and $T(n)$ of $P(n)$ and $Q(n)$, satisfy, at least approximately, the recurrences

$$S(n) = S(n-1) + 2T(n-1) + K_1,$$

and

$$T(n) = 2S(n-1) + 2T(n-1) + K_2,$$

for $n \in \mathbb{N}$, and with K_1, K_2 , nonnegative constants.

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FQMS and Grabiec Notions

Product of FQM Defined on the Domain of Words

Based on the notion of a fuzzy quasi-metric space [KM75].

Definition ([ChoGraRa06], [GregRom04])

A fuzzy quasi-metric on a set X is a pair $(M, *)$ such that $*$ is a continuous t-norm and M is a fuzzy set in $X \times X \times [0, \infty)$ such that for all $x, y, z \in X$:

$$(KM1) \quad M(x, y, 0) = 0.$$

$$(KM2) \quad x = y \text{ if and only if } M(x, y, t) = M(y, x, t) = 1 \text{ for all } t > 0.$$

$$(KM3) \quad M(x, z, t + s) \geq M(x, y, t) * M(y, z, s) \text{ for all } t, s \geq 0.$$

$$(KM4) \quad M(x, y, -) : [0, \infty) \rightarrow [0, 1] \text{ is left continuous.}$$

Definition ([ChoGraRa06], [GregRom04])

A fuzzy quasi-metric space is a triple $(X, M, *)$ such that X is a set and $(M, *)$ is a fuzzy quasi-metric on X .

FQMS and Grabiec Notions

Product of FQM Defined on the Domain of Words

Based on the notion of a fuzzy quasi-metric space [KM75].

Based on the notions of G-Cauchy sequence and G-complete fuzzy metric [Grab88] to obtain a fuzzy version of Banach fixed point theorem.

Definition

A sequence $(x_n)_n$ in a fuzzy metric space $(X, M, *)$ is **Cauchy** provided that $\lim_{n \rightarrow \infty} M(x_n, x_{n+p}, t) = 1$ for each $t > 0$ and $p \in \mathbb{N}$.

Definition

A fuzzy metric space $(X, M, *)$ is **complete** provided that every Cauchy sequence in X is convergent.

([SehBha72])

A **B-contraction** on a fuzzy metric space $(X, M, *)$ is a self-map f on X such that there is a constant $k \in (0, 1)$:

$$M(f(x), f(y), kt) \geq M(x, y, t) \text{ for all } x, y \in X, t > 0.$$

FQMS and Grabiec Notions

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Based on the notions of G-Cauchy sequence and G-complete fuzzy metric [Grab88] to obtain a fuzzy version of Banach fixed point theorem.

Thus, Grabiec's fixed point theorem can be formulated as follows.

Theorem ([Grab88])

*Let $(X, M, *)$ be a G-complete fuzzy metric space such that*

$\lim_{t \rightarrow \infty} M(x, y, t) = 1$ for all $x, y \in X$. Then every B-contraction on X has a unique fixed point.

Quasi-metric generalizations of B-contraction and G-completeness were introduced in [RoSaTi07] to define G-bicompleteness.

Theorem ([RoSaTi07])

*Let $(X, M, *)$ be a G-bicomplete fuzzy quasi-metric space such that*

$\lim_{t \rightarrow \infty} M(x, y, t) = 1$ for all $x, y \in X$. Then every B-contraction on X has a unique fixed point.

Non-Archimedean Bicompleteness

Product of FQM Defined on the Domain of Words

Although G -(bi)completeness is a very strong kind of completeness, we have the following nice fact for our approach.

Theorem ([RoSaTi07])

Each bicomplete non-Archimedean fuzzy quasi-metric space is G -bicomplete.

The following result implies that [RoSaTi07] Theorem applies to the standard fuzzy quasi-metric space of any bicomplete non-Archimedean quasi-metric space.

Theorem

Let (X, d) be a bicomplete non-Archimedean quasi-metric space. Then (X, M_d, \wedge) is a G -bicomplete (non-Archimedean) fuzzy quasi-metric space such that $\lim_{t \rightarrow \infty} M_d(x, y, t) = 1$ for all $x, y \in X$.

The Domain of Words

Product of FQM Defined on the Domain of Words

The domain of words Σ^∞ : all finite and infinite sequences (**words**) over an **alphabet** Σ , ordered by the so-called information order \sqsubseteq on Σ^∞ , i.e., $x \sqsubseteq y \Leftrightarrow x$ **is a prefix of** y .

Example

$$S \rightarrow aRR \rightarrow b|a|T|RT \rightarrow \lambda$$

a, aa, aab, aaba, aabab, ...

a, ab, abb, abba, ...

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Denote by $x \sqcap y$ **the longest common prefix** of x and y , and for each $x \in \Sigma^\infty$ denote by $\ell(x)$ the length of x .

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[Smyth88] introduced a non-Archimedean quasi-metric d_{\sqsubseteq} on Σ^∞ :

$$d_{\sqsubseteq}(x, y) = 0 \text{ if } x \sqsubseteq y, \text{ and } d_{\sqsubseteq}(x, y) = 2^{-\ell(x \sqcap y)} \text{ otherwise}$$

Its specialization order coincides with the order \sqsubseteq , and thus the quasi-metric space $(\Sigma^\infty, d_{\sqsubseteq})$ preserves the information provided by \sqsubseteq .

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Its specialization order coincides with the order \sqsubseteq , and thus the quasi-metric space $(\Sigma^\infty, d_{\sqsubseteq})$ preserves the information provided by \sqsubseteq .

Consequently d_{\sqsubseteq} is a bicomplete non-Archimedean quasi-metric on Σ^∞ .

Product FQMS of two FQMS

Product of FQM Defined on the Domain of Words

The **product (fuzzy quasi-metric) space** of two fuzzy quasi-metric spaces is the fuzzy quasi-metric space $(X_1 \times X_2, M_1 \times M_2, *)$ such that

$$(M_1 \times M_2)((x_1, x_2), (y_1, y_2), t) = M_1(x_1, y_1, t) * M_2(x_2, y_2, t).$$

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If (X_1, M_1, \wedge) and (X_2, M_2, \wedge) are non-Archimedean, then $(X_1 \times X_2, M_1 \times M_2, \wedge)$ is non-Archimedean.

*If $(X_1, M_1, *)$ and $(X_2, M_2, *)$ are bicomplete, then $(X_1 \times X_2, M_1 \times M_2, *)$ is bicomplete.*

Product FQMS of two FQMS

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The **product (fuzzy quasi-metric) space** of two fuzzy quasi-metric spaces is the fuzzy quasi-metric space $(X_1 \times X_2, M_1 \times M_2, *)$ such that

$$(M_1 \times M_2)((x_1, x_2), (y_1, y_2), t) = M_1(x_1, y_1, t) * M_2(x_2, y_2, t).$$

By applying the previous results and remarks to the standard fuzzy quasi-metric space of $(\Sigma^\infty, d_{\sqsubseteq})$ when $* = \wedge$, we deduce:

Theorem

$(\Sigma^\infty \times \Sigma^\infty, M_{d_{\sqsubseteq}} \times M_{d_{\sqsubseteq}}, \wedge)$ is a bicomplete non-Archimedean fuzzy quasi-metric space such that $\lim_{t \rightarrow \infty} (M_{d_{\sqsubseteq}} \times M_{d_{\sqsubseteq}})((x_1, x_2), (y_1, y_2), t) = 1$.

Therefore, every **B-contraction** on this space has a **unique fixed point**.

Functional

Product of FQM Defined on the Domain of Words

We build a functional Φ suggested by S and T .

$$\Phi : \Sigma^\infty \times \Sigma^\infty \rightarrow \Sigma^\infty \times \Sigma^\infty,$$

given for each pair $x^1, x^2 \in \Sigma^\infty$, by $\Phi(x^1, x^2) = (u^1, u^2)$,
where $(u^1)_0 = S(0)$, $(u^2)_0 = T(0)$, and

$$(u^1)_n = x_{n-1}^1 + 2(x^2)_{n-1} + K_1,$$

$$(u^2)_n = 2(x^1)_{n-1} + 2(x^2)_{n-1} + K_2,$$

for all $n \in \mathbb{N}$ such that $n \leq (\ell(x^1) \wedge \ell(x^2)) + 1$.

Functional

Product of FQM Defined on the Domain of Words

We build a functional Φ suggested by S and T .

We show that Φ is a B-contraction of the G-bicomplete
(non-Archimedean) fuzzy quasi-metric space $(\Sigma^\infty \times \Sigma^\infty, M_{d_{\sqsubseteq}} \times M_{d_{\sqsubseteq}}, \wedge)$.

Hence Φ **has a unique fixed point which is the solution of the recurrences S and T .**

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To model the algorithm iterations, one actually works on the set of all finite words, that endowed with the restriction of $(M_{d_{\sqsubseteq}}, \wedge)$ provides a non-bicomplete non-Archimedean fuzzy quasi-metric space.

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Hence Φ **has a unique fixed point which is the solution of the recurrences S and T .**

The product space $(\Sigma^F \times \Sigma^F, M_{d_{\sqsubseteq}} \times M_{d_{\sqsubseteq}}, \wedge)$ is also a non-bicomplete non-Archimedean fuzzy quasi-metric space. It is bicompletable and its bicompletion is a non-Archimedean fuzzy quasi-metric space.

In particular, for each pair x^1, x^2 in the set of finite words, the sequence of iterations converges, in $(\Sigma^\infty \times \Sigma^\infty, (M_{d_{\sqsubseteq}} \times M_{d_{\sqsubseteq}})^i, \wedge)$, to the solution for the pair of recurrence equations S and T .

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Bicompletion of QMS

Product QMS of Complexity Space

Definition

The product quasi-metric space of two quasi-metric spaces (X, d) and (Y, e) is the quasi-metric space $(X \times Y, d \times e)$, where $d \times e$ is

$$(d \times e)((x_1, y_1), (x_2, y_2)) = d(x_1, x_2) \vee e(y_1, y_2),$$

for all $(x_1, y_1), (x_2, y_2) \in X \times Y$.

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for all $(x_1, y_1), (x_2, y_2) \in X \times Y$.

Definition

Given a quasi-metric d on X , then the function d^{-1} defined by $d^{-1}(x, y) = d(y, x)$, is also a quasi-metric on X , called the **conjugate** of d , and the function d^s defined by $d^s(x, y) = d(x, y) \vee d^{-1}(x, y)$ is a metric on X .

Definition

A quasi-metric space (X, d) is said to be **bicomplete** if (X, d^s) is a complete metric space.

Contraction

Product QMS of Complexity Space

By a **contraction map** on a quasi-metric space (X, d) we mean a self-map f of X such that $d(fx, fy) \leq kd(x, y)$ for all $x, y \in X$, where k is a constant with $0 < k < 1$. The number k is called a contraction constant for f .

The classical Banach contraction principle can be generalized to the quasi-metric setting as follows (see for instance Lemma 2.4 of [1]).

Theorem

Let f be a contraction map on a bicomplete quasi-metric space (X, d) . Then, for each $x \in X$, the sequence of iterations $(f^n x)_{n \in \omega}$ is convergent in (X, d^s) to a point $x_0 \in X$ which is the unique fixed point of f .

Schellekens Complexity Space

Product QMS of Complexity Space

[Sch95] **complexity (quasi-metric) space**, $(\mathcal{C}, d_{\mathcal{C}})$

$$\mathcal{C} = \left\{ f : \omega \rightarrow (0, \infty] : \sum_{n=0}^{\infty} 2^{-n} \frac{1}{f(n)} < \infty \right\},$$

and $d_{\mathcal{C}}$ is the quasi-metric on \mathcal{C} given by

$$d_{\mathcal{C}}(f, g) = \sum_{n=0}^{\infty} 2^{-n} \left(\left(\frac{1}{f(n)} - \frac{1}{g(n)} \right) \vee 0 \right)$$

for all $f, g \in \mathcal{C}$.

Schellekens Complexity Space

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for all $f, g \in \mathcal{C}$.

Consequence of Theorem 1, and Remark on p. 317 [RomSch99]:

Theorem

The complexity space $(\mathcal{C}, d_{\mathcal{C}})$ is bicomplete.

Functional

Product QMS of Complexity Space

We construct a monotone increasing functional Φ , which is a contraction on $(\mathcal{C} \times \mathcal{C}, d_{\mathcal{C}} \times d_{\mathcal{C}})$.

Theorem

Let Φ be the functional on $\mathcal{C} \times \mathcal{C}$ defined by

$$\Phi(f, g)(0) = (S(0), T(0)),$$

and

$$\begin{aligned}\Phi(f, g)(n) = & (pf(n-1) + qg(n-1) + K_1, \\ & rf(n-1) + sg(n-1) + K_2),\end{aligned}$$

for $n \in \mathbb{N}$ and $f, g \in \mathcal{C}$.

(cont...)

Functional

Product QMS of Complexity Space

We construct a monotone increasing functional Φ , which is a contraction on $(\mathcal{C} \times \mathcal{C}, d_{\mathcal{C}} \times d_{\mathcal{C}})$.

Theorem

(... cont)

If $\alpha < 1$, where

$$\alpha = \frac{1}{2} \left(\frac{1}{p \wedge r} + \frac{1}{q \wedge s} \right),$$

then:

- (1) Φ is a monotone increasing contraction on $(\mathcal{C} \times \mathcal{C}, d_{\mathcal{C}} \times d_{\mathcal{C}})$ with contraction constant α .
- (2) Φ has a **unique fixed point** (f_0, g_0) .

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Preorder

Preordered Complete FQMS

Take as a starting point Ricarte and Romaguera Theorem 2.2 [RicRom13] new version of Matowski's theorem (based on the notion of a t -norm of Hadzić), which generalized Banach's contraction principle using a type of contractions that generalize Hicks' C-contractions.

Theorem ([RicRom13])

. Let $(X, M, *)$ be a complete fuzzy metric space and $f : X \rightarrow X$ a self-map such that

$$M(x, t, y) > 1 - t \rightarrow M(fx, fy, \varphi(t)) > 1 - \varphi(t),$$

for all $x, y \in X$ and $t > 0$, where $\varphi : [0, \infty) \rightarrow [0, \infty)$ is a nondecreasing function satisfying $\lim_{n \rightarrow \infty} \varphi^n(t) = 0$ for all $t > 0$. Then f has a unique fixed point

Preorder

Preordered Complete FQMS

A preorder on a nonempty set X is a reflexive and transitive binary relation \preceq on X . A preordered fuzzy (quasi-)metric space is a 4-tuple $(X, M, \preceq, *)$ such that $(X, M, *)$ is a fuzzy (quasi-)metric space and \preceq is a preorder on X .

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If $(M, *)$ is a fuzzy quasi-metric on X , the relation \leq_M on X given by $x \leq_M y \Leftrightarrow M(x, y, t) = 1$ for all $t > 0$, is an order on X called the specialization order of $(M, *)$.

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Theorem ([CasRomTi14])

*If the ordered fuzzy quasi-metric space $(X, M, \leq_M, *)$ is \leq_M -complete and $f : X \rightarrow X$ is a \leq_M -nondecreasing self-map such that there is $x_0 \in X$ satisfying $x_0 \leq_M f x_0$, then f has a fixed point.*

Complexity Space

Preordered Complete FQMS

To deduce the existence of the solution for the algorithm equations we shall use the following subset of the complexity space from [Sch95]:

$$\mathcal{C}_1 = \{f \in \mathcal{C} : f(n) \geq 1 \text{ for all } n \in \mathbb{N}\}$$

and the function $Q_{\mathcal{C}}$ defined in [RomTi09]:

$$Q_{\mathcal{C}}(f, g, t) = \sum_{k=n}^{\infty} 2^{-k} \left(\left(\frac{1}{g(k)} - \frac{1}{f(k)} \right) \vee 0 \right),$$

where $t \in (n-1, n]$, $n \in \mathbb{N}$.

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where $t \in (n-1, n]$, $n \in \mathbb{N}$.

Now we can construct a fuzzy set M_1 in $(\mathcal{C} \times \mathcal{C}) \times (\mathcal{C} \times \mathcal{C}) \times [0, \infty)$ as:

$$M_1((f_1, g_1), (f_2, g_2), 0) = 0 \quad \text{for all } f_i, g_i \in \mathcal{C}_1,$$

$$M_1((f_1, g_1), (f_2, g_2), t) = 1 \quad \text{if } f_1 \leq f_2 \text{ and } g_1 \leq g_2 \text{ and } t > 0, \quad \text{and}$$

$$M_1((f_1, g_1), (f_2, g_2), t) = 1 - [Q_{\mathcal{C}}(f_1, f_2, t) \vee Q_{\mathcal{C}}(g_1, g_2, t)] \quad \text{otherwise.}$$

Fuzzy Quasi-Metric Space

Preordered Complete FQMS

This set must be appropriate to define the following fuzzy quasi-metric space:

Lemma (Lemma 1 [CasRomTi14])

$(\mathcal{C}_1, M_1, \leq_{M_1}, *_L)$ is a \leq_{M_1} -complete fuzzy quasi-metric space.

So we will have to show that:

$$Q_C(f_1, f_2, t + s) \leq Q_C(f_1, f_3, t) + Q_C(f_3, f_2, s)$$

$$Q_C(g_1, g_2, t + s) \leq Q_C(g_1, g_3, t) + Q_C(g_3, g_2, s)$$

and that the fuzzy quasi-metric space is \leq_{M_1} -complete.

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$$Q_C(g_1, g_2, t + s) \leq Q_C(g_1, g_3, t) + Q_C(g_3, g_2, s)$$

and that the fuzzy quasi-metric space is \leq_{M_1} -complete.

We have the following consequence Theorem 6 of [CasRomTi14]

Theorem ([CasRomTi14])

If $\Phi : \mathcal{C}_1 \rightarrow \mathcal{C}_1$ is a \leq_{M_1} -nondecreasing map and there is $f_0 \in \mathcal{C}_1$ such that $f_0 \leq_M \Phi f_0$, then Φ has a fixed point.

Contraction Condition

Preordered Complete FQMS

One of the most interesting facts of this new approach is that the contraction condition of the preceding theorem is automatically satisfied whenever the self-map f is nondecreasing for the specialization order and $\varphi : [0, \infty) \rightarrow [0, \infty)$ verifies that $\varphi(t) > 0$ for all $t > 0$.

We will have to show that if $f_1 \leq f_2$ and $g_1 \leq g_2$ then $\Phi(f_1, g_1) \leq \Phi(f_2, g_2)$ and then by the preceding Theorem we will find that Φ has a fixed point, which is the solution to the system of recurrence equations.

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Introduction

Methods to Study the Existence and Uniqueness of Solution

Review previous approaches and sketch the application of a new approach [CasRomTi14]:

- Product of fuzzy quasi-metric defined on the domain of words [CasRomTi10b].
- Product quasi-metric space of complexity space [CasRomTi10a].
- Preordered complete fuzzy quasi-metric space.